

DAG-Edit & OBO

- Come on a magical journey of discovery with me, and learn all the new things DAG-Edit can do now that we have the OBO format...

The Wonderful World of Relationship Types

- OBO has imbued relationship types with all kinds of new and amazing powers.

Range

- Relationship types now have a “range” attribute.
 - The range is a term. Only terms that are subclasses of the range term can be the target of a term relationship type.

Range

- So, let's say we have the relationship type *has_sequence*. The range of *has_sequence* is *sequence*. That means that a relationship with the type *has_sequence* can only have parents that are subclasses of the term *sequence* (like *rna_sequence*, *dna_sequence*, etc).

Legal:



Illegal:



Range

- This gives us a lot more ability to make sure the ontologies are consistent. This becomes really important when you're dealing with an ontology with a lot of relationship types (like the Sequence Ontology, or an anatomy, or the Phenotype Ontology, or any ontology with “cross products”).

Domain

- Relationship types also can specify a domain.
 - The domain is a term. Only terms that are subclasses of the domain term can be the source of a term relationship type.

Domain

- So, let's say we have the relationship type *has_color*. The domain of *has_color* is *visible_object*. That means that a relationship with the type *has_color* can only have children that are subclasses of the term *visible_object* (like *chicken*, *rock*, *leonardo_da_vinci*, etc).

Legal:



Illegal:



Is Cyclic

- It is possible to say whether or not a relationship type can be used to create a cycle. It is illegal to create a cycle of relationships out of a relationship type that is not marked “is cyclic”.
- However, it is still legal to create cycles using multiple acyclic relationship types in combination.

Is Transitive

- Transitive relationship types are relationship types such that:
 - If A has relationship X to B, and B has relationship X to C, by definition, A has relationship X to C.
 - is_a and part_of are transitive, which is why the true path rule works.

Is Symmetric

- Symmetric relationship types are relationship types such that:
 - If A has relationship X to B, B has relationship X to A.
 - Precise identity relationships (like mathematical equality) are symmetric (if $A=B$, $B=A$). `is_a` and `part_of` are not symmetric.

Namespaces

- It is now possible to assign terms (and term relationships) to a “namespace”.
 - Other things, like categories and dbxrefs, will be assignable to namespaces someday.

Namespaces

- A namespace can be thought of in two ways
 - A namespace refers to the logical ontology to which a term belongs, regardless of where the ontology is stored. Namespaces let us know which ontology obsolete terms belong to, since we don't have any parentage info.
 - A namespace refers to the file in which a term (or relationship, or whatever) should be stored. The OBO save mechanism lets you save your session into several files based on namespaces.

Namespaces

- Why bother with namespaces instead of assigning file names?
 - Terms might have been loaded from a source that doesn't have a proper name, like a database, so we use a more abstract designation
 - Terms in different files might still belong to the same namespace (if, for example, you created a sub-ontology in a special file, and wanted to merge it with another ontology).

Relationship Namespaces

- A relationship can be marked with a different namespace than it's parent or child.
 - This can be used to create a distinct file for relationship types that don't belong in an existing ontology. For example, you could a bunch of links between two ontologies, you could keep those links in a third file by giving them a different namespace than the other two ontologies.
 - Relationships with no namespace are assigned to the namespace of the child term.

Inverse Necessity

- It is possible to mark a relationship as necessarily true.
 - A necessarily true relationship *must always* be true. (ie *human_finger* is necessarily part_of *human_hand*, but *shoe* is not necessarily part_of *casual_outfit*, because a shoe is sometimes part_of a casual outfit, but someone might go barefoot).

Inverse Necessity

- It is possible to mark a the inverse of a relationship as necessarily true. A relationship may be inverse necessarily true with being necessarily true
 - For example: The relationship *wheel part_of car* is inverse necessarily true. Something with no wheels is not a car, but a wheel can still be a wheel without being part of a car.

Completeness

- A “complete” term specification is one where anything that matches the term specification is *by definition* an instance of that term.

Example: An Incomplete Spec

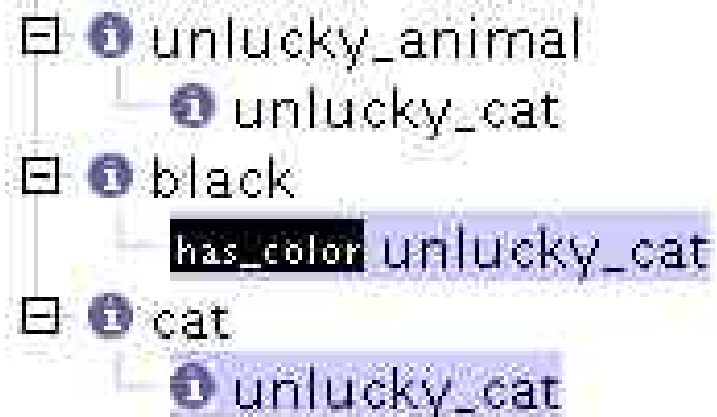
- The following term specification of `manx_cat` is incomplete:



- Because being a `tailless_animal` and being a `cat` does not automatically make something a `manx_cat` (it might be `cat_with_severed_tail`).

Example: An Complete Spec

- The following term specification of `unlucky_cat` is complete:



- The relationships **has_color *black*** and **is_a *cat*** are enough to constitute a complete definition, because anything that is black and a cat is an unlucky cat.
- A complete term spec can contain relationships that don't contribute to the completeness (like **is_a *unlucky_animal***).

The Parent Plugin

- There's a single easy place to set all these relationship attributes: The Parent Plugin.

Fun with Obsolete Terms

- It is now possible to create a link between an obsolete term and the terms that should be used to replace it.

New Built-In Relationship Types

- `disjoint_from`
- `inverse_of`

disjoint_from

- A is *disjoint_from* B if no instance of A can ever be an instance of B.
 - *smart_guy* is disjoint_from *dumb_guy*
 - *visible_object* is disjoint_from *abstract_concept*
 - disjoint_from applies to subclasses of the specified terms. So if *george_w* is_a *dumb_guy*, *smart_guy* is disjoint_from *george_w*

inverse_of

- `inverse_of` applies only to relationship types. It's easier to explain with examples:
 - *part_of* is the `inverse_of` *has_part*
 - *is_a* is the `inverse_of` *has_subclass*
 - *has_color* is the `inverse_of` *is_color_of*
- `inverse_of` is symmetrical. If A is the inverse of B, B is the `inverse_of` A.

Conclusion: How does this help us

- All of these OBO-centric concepts are already present in the GO, they just aren't explicitly encoded. Hence the ongoing discussions about whether `part_of` means necessarily part of; sometimes it does, sometimes it doesn't, we just had no way of making it clear).
- With this additional information, we can now map to OWL and interact with reasoners much more effectively.